## Site Reliability Engineering Distributed PubSub

Non-Abstract Large System Design



#### NALSD

- "Non-Abstract Large System Design"
- Alternatively: SRE Classroom
- Large ("planet scale") system design questions
- Hands-on workshops and exercises
- Non-abstract component:
  - Crunch numbers
  - Provision the system
- Resilient software systems
- Distributed architecture patterns

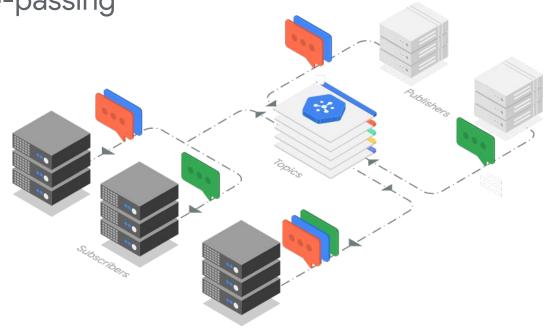
#### Agenda

- Introduction and problem statement
- "Let's do it together"
- Breakout session 1: **Design for single datacenter**
- Single datacenter sample solution
- Breakout session 2: **Design for multiple datacenters**
- Multiple datacenters sample solution
- Breakout session 3: Provision the system
- Provision the system sample solution
- Wrap-up and conclusions

#### Introduction



- Publish-Subscribe (PubSub)
- Asynchronous communication through message-passing



- Publishers: "producers" or "writers"
  - Senders of messages
  - Sends ordered messages
  - Messages grouped by topic

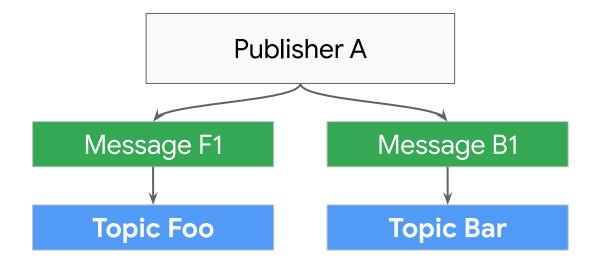
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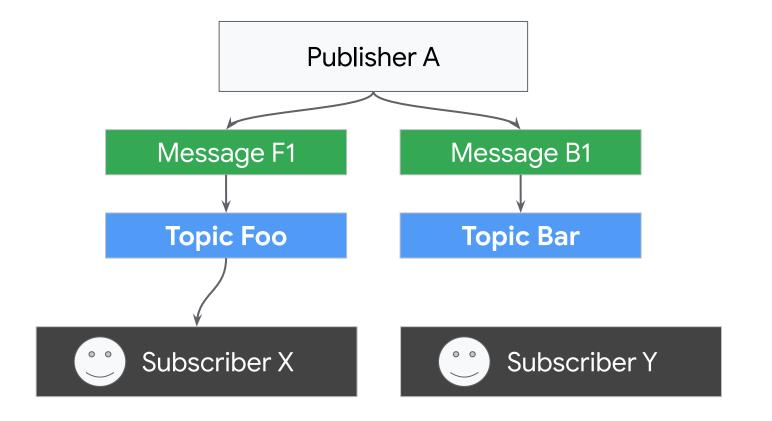
- Subscribers: "consumers" or "readers"
  - Subscribes to topics
  - Receives messages only for subscribed topics

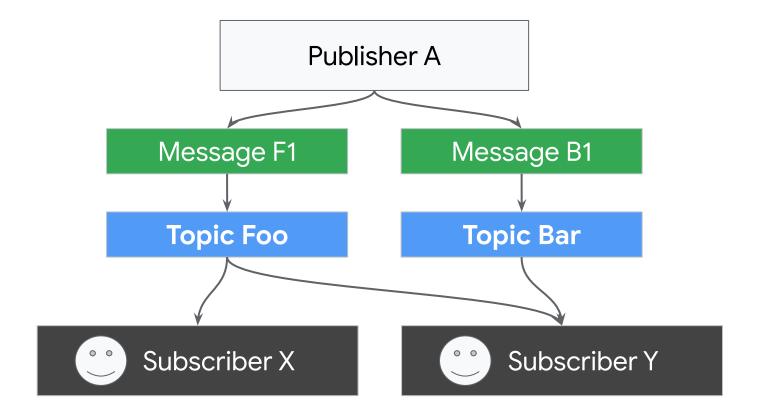


- Publishers do not directly communicate with Subscribers
- Subscribers **do not directly communicate** with Publishers
  - Scale publishers/subscribers independently SLIDSCITICOIS









#### Problem Statement Let's identify the problem at hand



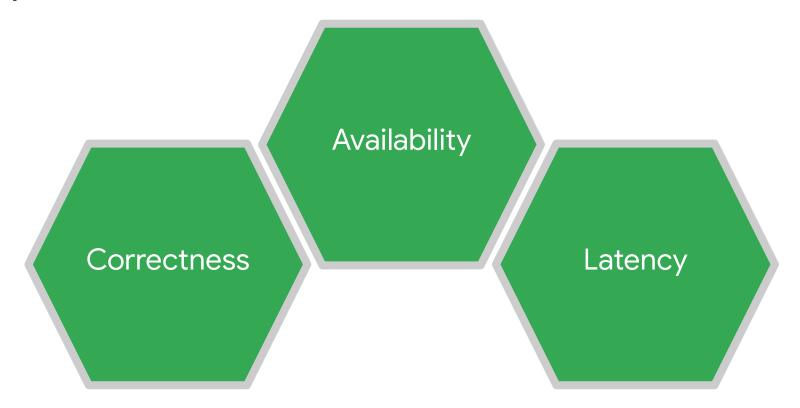


#### Design a PubSub service that clients all over the world can use to read and write messages.

#### Gather Requirements Let's identify what we know and what we need



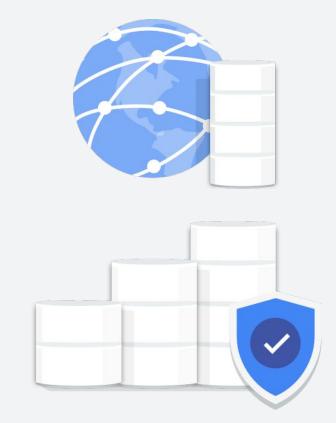
#### Requirements



#### Background

What we have:

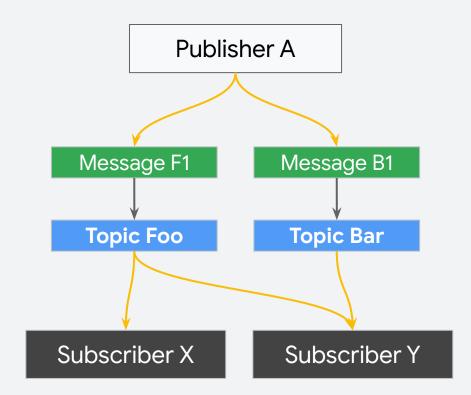
- Three datacenters (DCs):
  - New York
  - Seattle
  - Kansas City
- Reliable storage system
   Distributed!
- Reliable network
- Authentication & Authorization



#### Requirements

What we need:

- A way to publish messages
  - Ordered
  - Grouped by topic
- A way to receive messages
  - Ordered
  - Grouped by topic
- Message persistence



#### Requirements

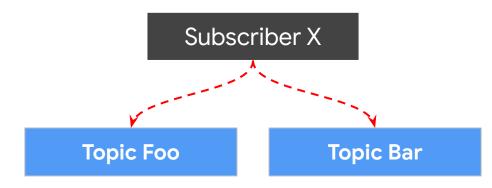
- Each DC runs the PubSub service we are designing
- Clients all over the world read and write messages
- Large volume of messages per day
- Uneven distribution of traffic over time



#### Requirements - What Does PubSub Do?

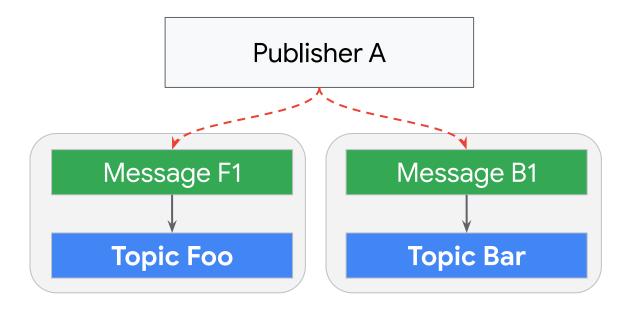
- Communicate ordered messages, grouped by topic
- Readers/writers can connect to any DC
- Users expect the same level of service from all DCs
- If a DC goes down, the user will automatically get connected to another one (this is already provided as a service)
- Once a DC recovers, it goes back to full service

- Topics are identified by their **topic\_id**.
- Readers are identified by their **consumer\_id**.
- Readers will explicitly subscribe to topics.
- Subscribe(topic\_id, consumer\_id): Subscribes the given consumer to the given topic.

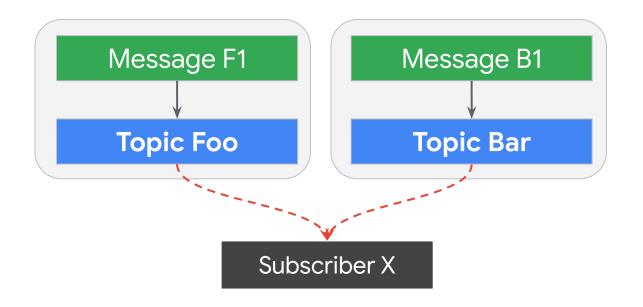


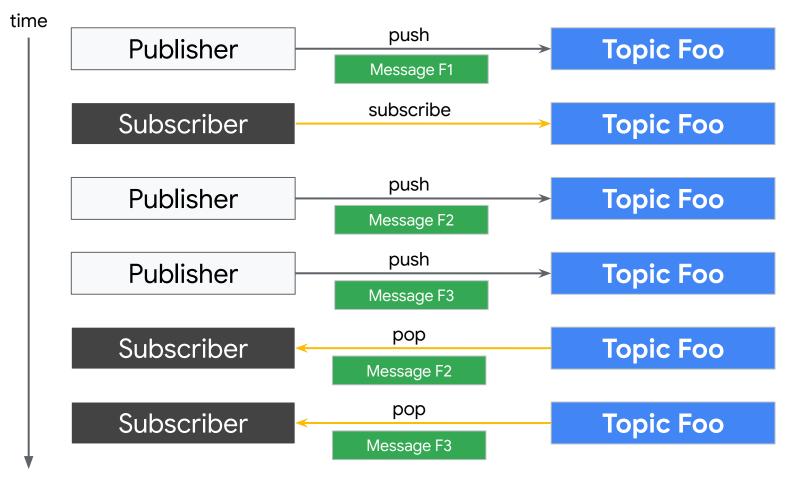
• Push(topic\_id, message):

Append the message to the given topic.



• **Pop(topic\_id, consumer\_id):** Read the next message (in order) for the given topic.





• List():

Returns a list of all available topics.

• Not in scope for this exercise.



#### Service Level Terminology

• SLI: service level indicator

A quantifiable (numeric) measure of service reliability.

• SLO: service level <u>objective</u>

A reliability target for an SLI.

• SLA: service level <u>agreement</u>

SLO + consequences when SLO is violated

#### Requirements - SLO Availability

 PubSub must continue working under peak load even if one datacenter goes down

#### Latency

- 99% of API calls must complete within 500ms
- 99% of pushed messages must be available for pop anywhere in the world within 1s

#### Requirements - SLO

#### Correctness

- At-Least-Once delivery
- 100 day message retention
- System can lose 0.01% of enqueued message per year

## Further details, including volumes of data, are in the workbook handouts.

### Let's do it together: push()



#### **Requirements Recap**

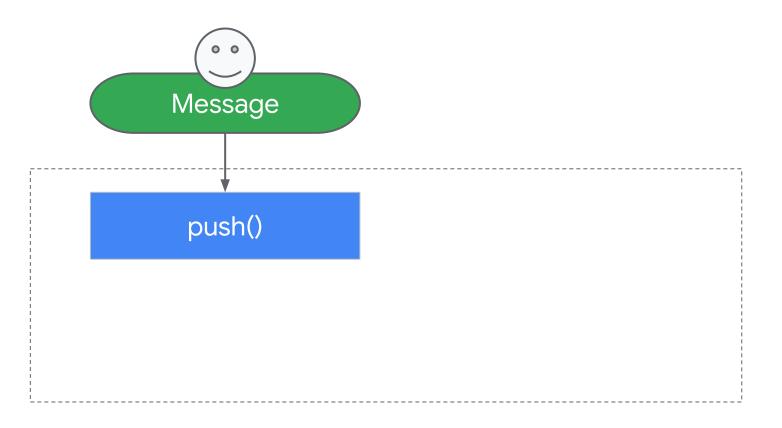
- Global PubSub Service
- Three datacenters (DCs):
  - New York
  - Seattle
  - Kansas City
- Clients all over the world write (push) and read (pop)
- Large volume of messages per day
- Uneven distribution of traffic over time

### push()

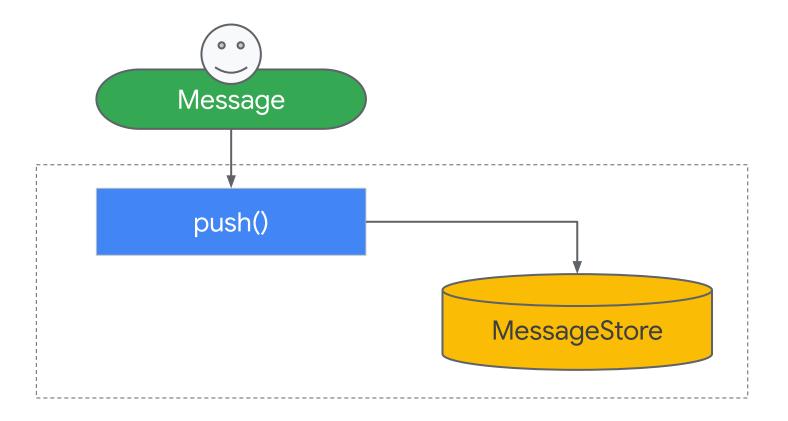
# Let's design the API call that receives messages.



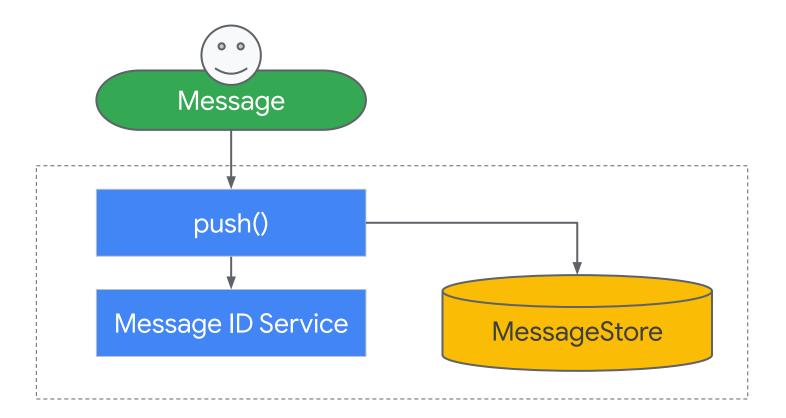
#### Pushing a message



#### Start by storing the messages...

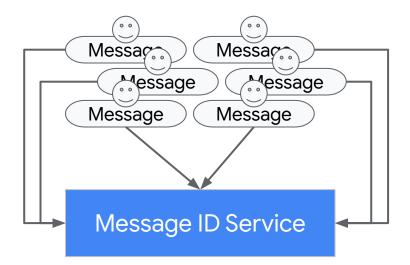


#### Assign message IDs for storage...



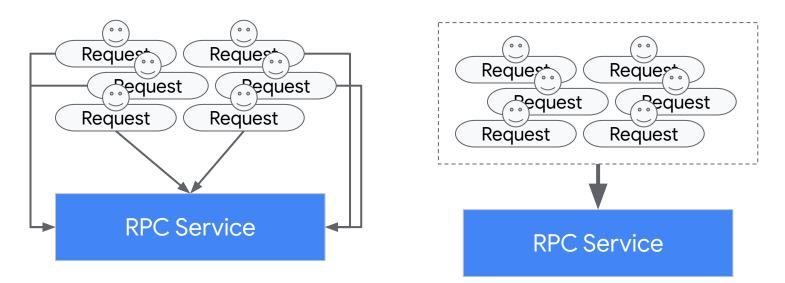
#### More on the Message ID Service

- Assign **unique** IDs for message within a topic
- Assign ordered message IDs for simple ordered lookup



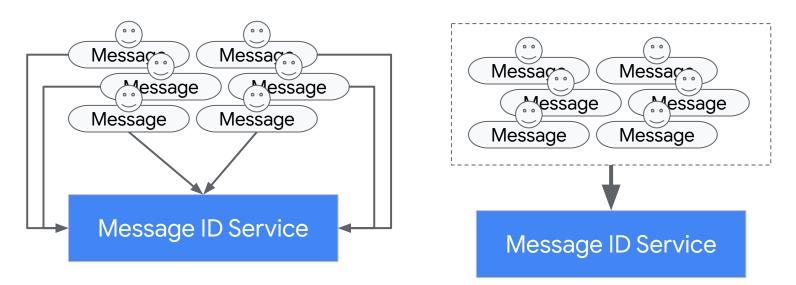
#### **Batch Operations**

- Address **bandwidth** or **throughput** bottlenecks
- May be supported alongside singular operations
- Basically: stuff multiple requests into a single RPC



#### More on the Message ID Service

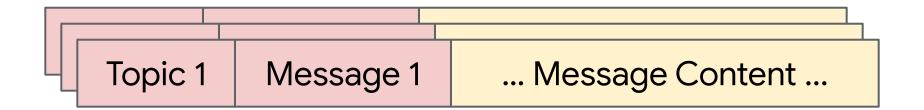
- Assign **unique** IDs for message within a topic
- Assign ordered message IDs for simple ordered lookup
- Performance optimizations: batch operations



### More on the MessageStore



#### Key: Topic ID, Message ID Value: Message Content



# More on the MessageStore

- Distributed file system
  - Storage abstractions
  - write(), read(), implemented already
  - Supports configurable replication strategy

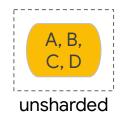


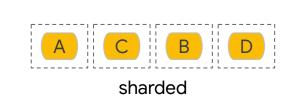
# Message Store Sharding

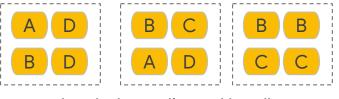
- Need to retain 100 days worth of messages
- 100 days \* ... = 25TB of data  $\rightarrow$  too big for one machine :(

# Sharding

- Address **storage size** bottlenecks
- Basically: split your data into multiple buckets, and store those buckets separately, possibly multiple copies of each bucket
- Sharding mechanism should be flexible
- Consistency and fault tolerance
- A single disk failure should not cause data loss
- Consider replicating shards locally (local reads are cheapest)





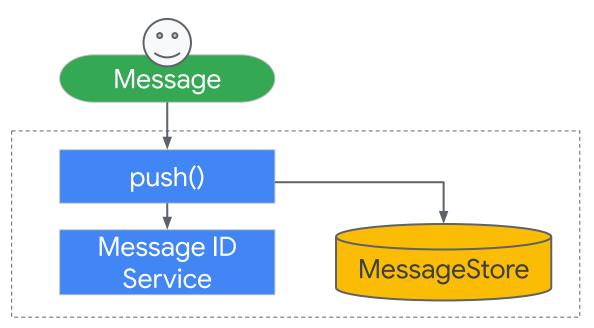


# Message Store Sharding

- Need to retain 100 days worth of messages
- 100 days \* ... = 25TB of data  $\rightarrow$  too big for one machine :(
- Sharding to the rescue!
- Keep multiple copies (replicas) of each shard:
  - Greater resilience
  - ... and performance too (local reads are cheap)!

# Flow overview: push()

- 1. Get message ID from Message ID Service
- 2. Write message to MessageStore
- 3. Ack receipt of message



# Reminder: don't sweat it!

- Designs will be different, with different abstractions: that's okay!
- Focus on the process of designing something end-to-end
- Think about high level concepts, rather than nitty details
- Think about trade-offs of different design decisions
- Make assumptions explicit
- Call out risks
- Simplify the problem
- If working in a group, discuss ideas and use each other as resources!

# Rules of engagement

- Assume good intent
- Respect each other
- Speak up and share information
- Let everybody speak
- Ask questions

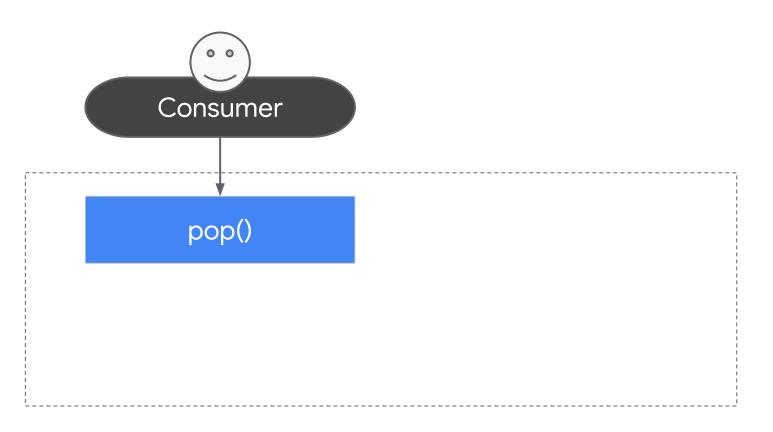
#### Most importantly, have fun!

# Breakout Session 1: Single Datacenter (40 minutes)

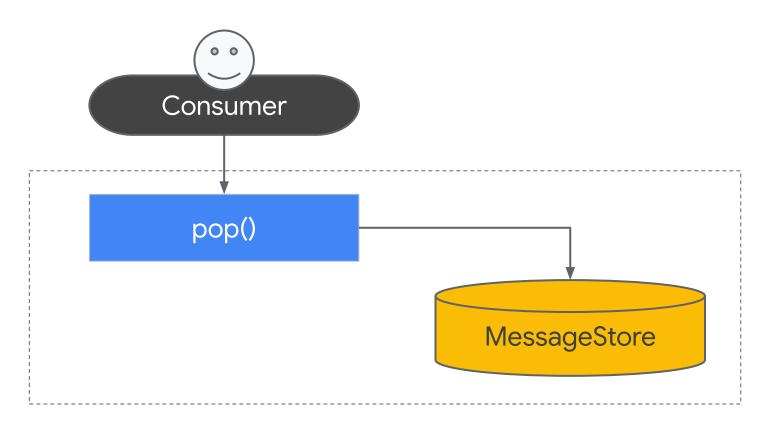
# Goal: Design a working system that fits in a single datacenter.

# **Break: 5 Minutes**

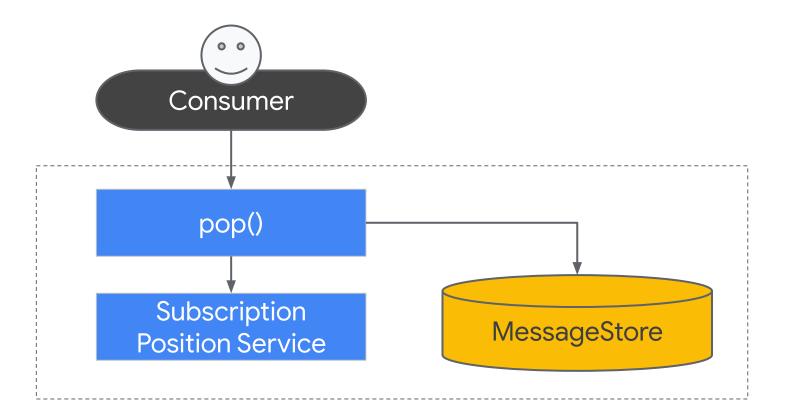
# Reading a message



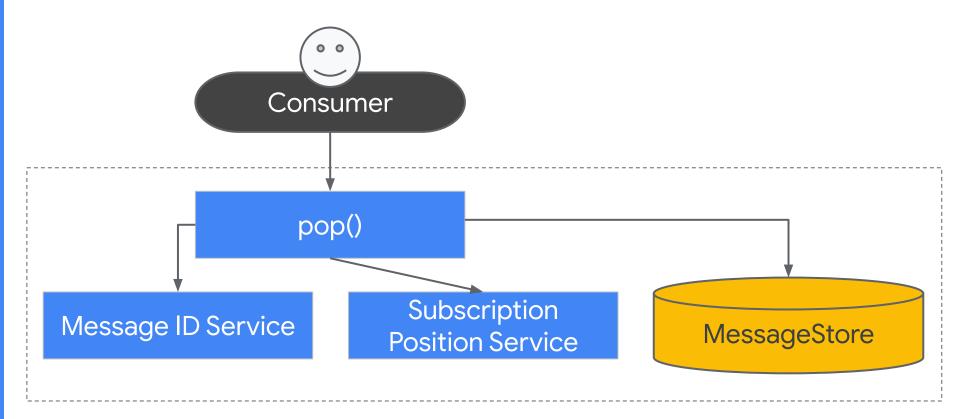
# Reading a message



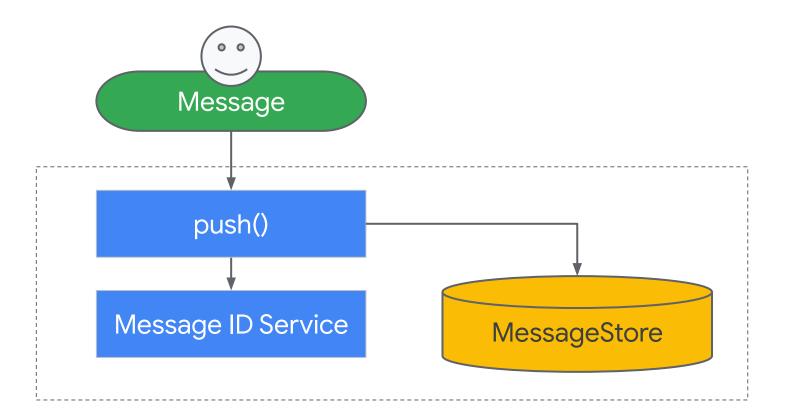
### Reading: getting the "next" message



### Next, read the messages on demand...

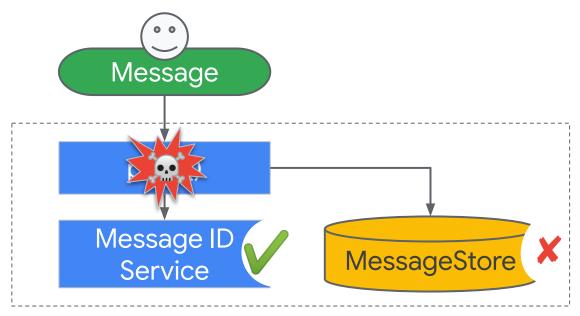


### Reminder of how push() works...



# Error Handling: pop()

- Message IDs are consecutive... almost.
- Gaps can arise if push() service crashes after allocating ID, but before message is successfully written to storage.

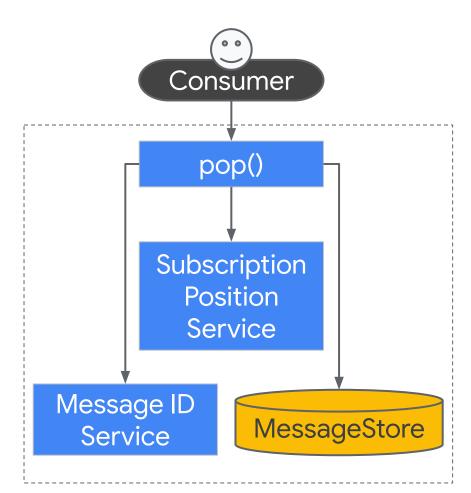


# Error Handling: pop()

- Detect error upon read
- Increment ID and keep reading until the next message is found
- Do not read past the end of the topic
- Some latency impact; expect to be rare
- Performance optimizations:
  - Batch reads
  - Readahead cache
  - Bloom filter on storage service

# Flow Overview: pop()

- 1. Get latest written message ID from Message ID Service
- 2. Get latest read message ID from Subscription Position Service
- 3. Increment the read message ID
- 4. If at the end of topic, return
- 5. Read message from storage
- 6. Return the message to consumer
- 7. Update subscription position for consumer and topic

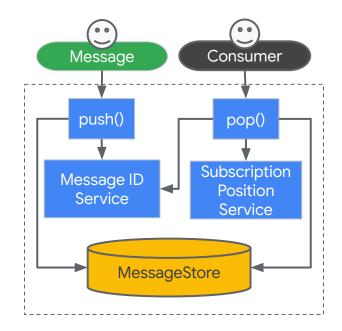


# Breakout Session 2: Multiple Datacenters (30 minutes)

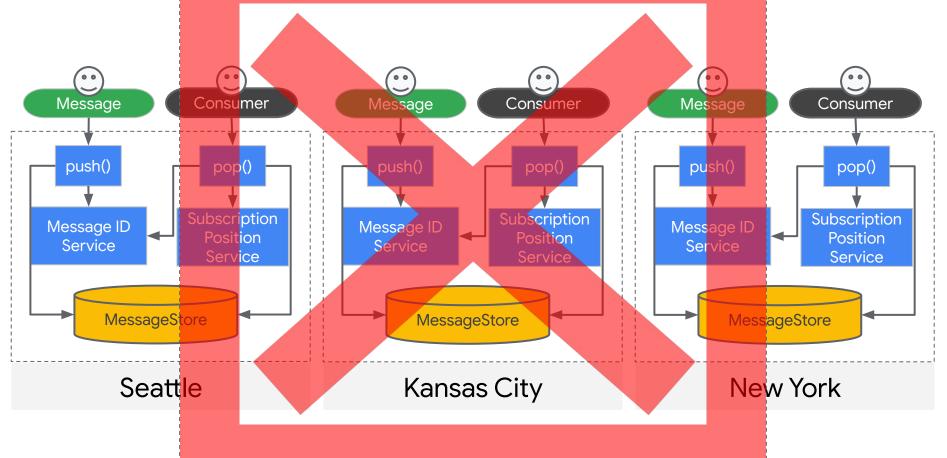
# Goal: Extend the design to work correctly in multiple datacenters.

# **Break: 5 Minutes**

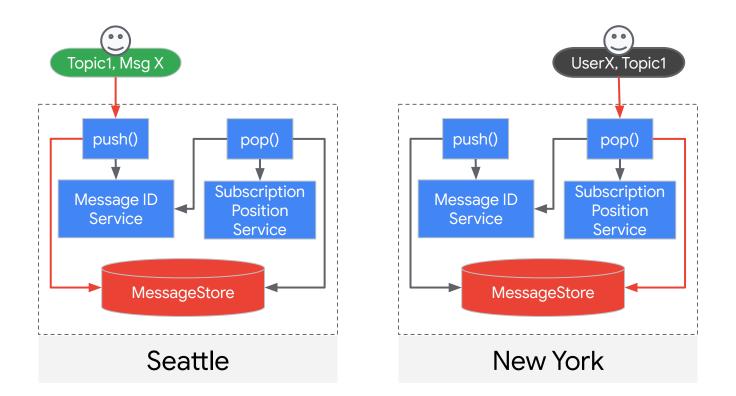
# Single Datacenter Design



### One for each datacenter...?



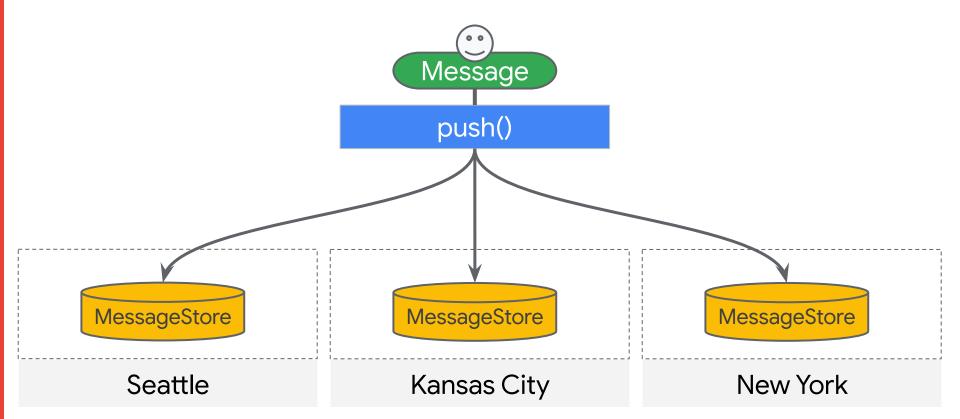
### Partitioned MessageStore



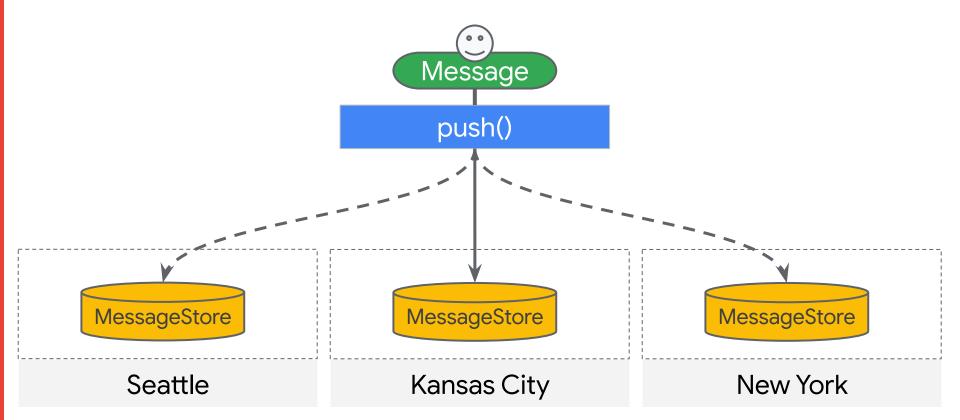
# MessageStore Replication

- Pushes can arrive at any datacenter
- Need to be able to pop messages from any datacenter, even at a different datacenter than where it arrived
- Need to replicate messages to every datacenter
- Factors to consider:
  - Consistency
  - Fault tolerance
  - Availability

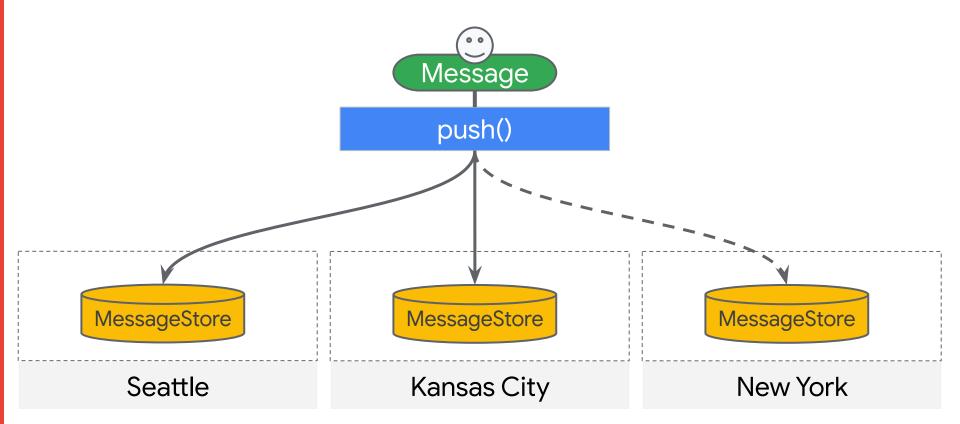
# Replication: synchronous



# Replication: asynchronous



### **Replication:** hybrid



# MessageStore Replication: Tradeoffs

	Push Latency	Pop Latency	Data Durability
Synchronous Replication	High	Low	High
Asynchronous Replication	Low	High	Low
Hybrid Replication	Medium	Medium	Medium

# MessageStore Replication

- Asynchronous writes: ~10ms response time
- Can we afford the data loss?
- Reminder:
  - Can lose 0.01% of pushed messages per year
  - 99% of messages must be available for pop from any location in 1 second or less

5,000 topics \* 10,000 msg / day / topic = 50M msg / day  $\rightarrow$  Can lose 5k messages per day.

# Async Replication

90k sec/day \* 1 msg/sec/thread

= 90k msg / day / thread

parallelize processing to handle the entire load...

(50M msg / day) /

(90k msg / thread) =

~600 threads / day

(i.e. **concurrent** loads / day)

Reminders:

- 50M msg / day
- 99% of messages must be available for pop from any location in 1 second or less
- ~90k seconds / day
- Assume 1 second replication delay

# Async Replication

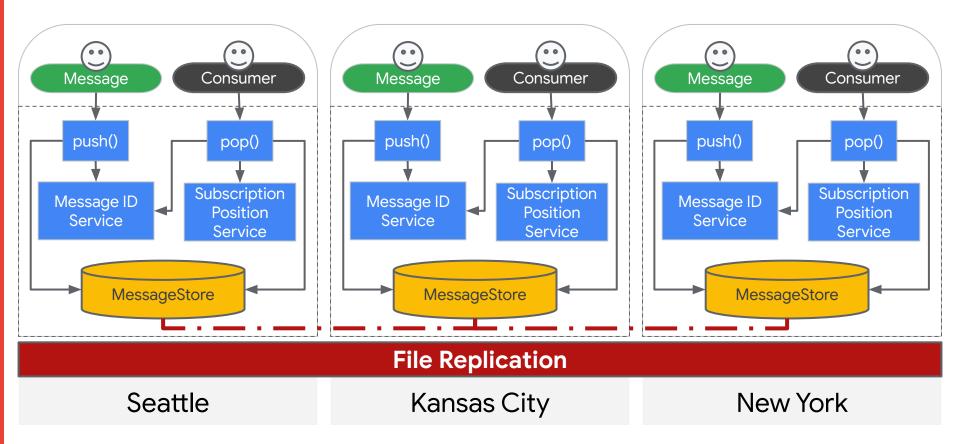
- Machine would have to fail ~8 times / day for us to lose 5k messages (0.01% of incoming messages)

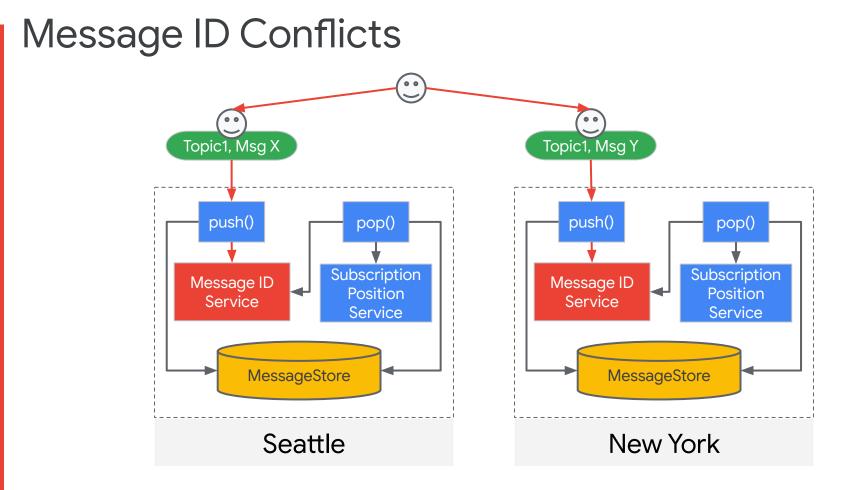
We can afford it!

Reminders:

- Can lose 5k msg / day
- ~600 in-flight msg / sec

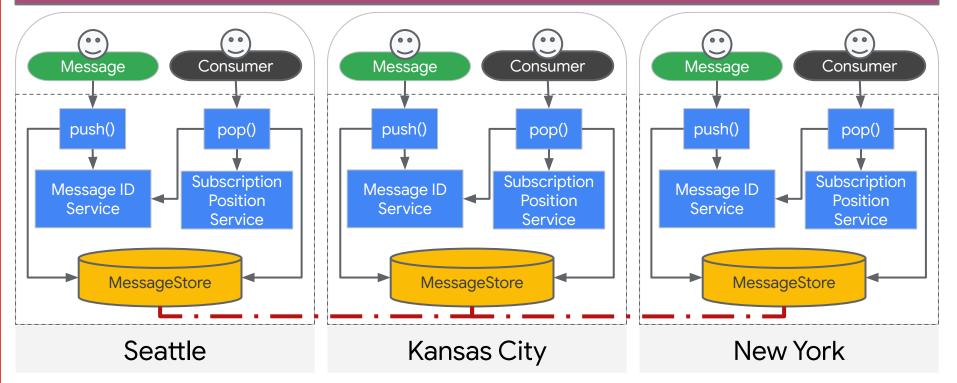
### Let's use replication...





### Let's use consensus...

#### **Paxos-based consensus**

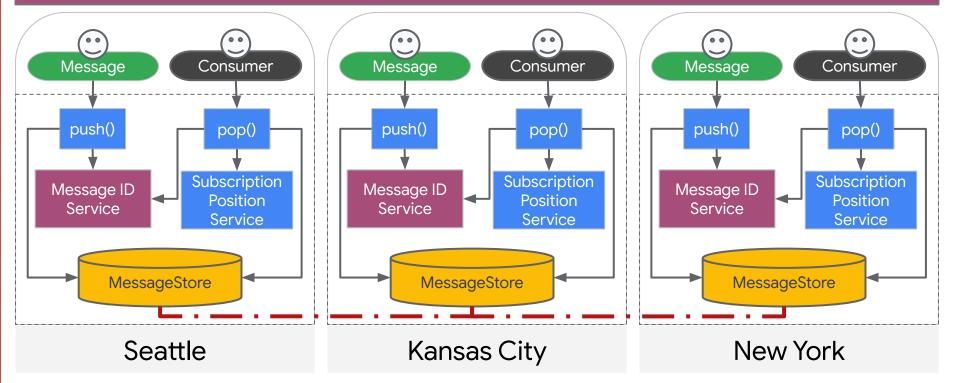


# **Distributed Consensus**

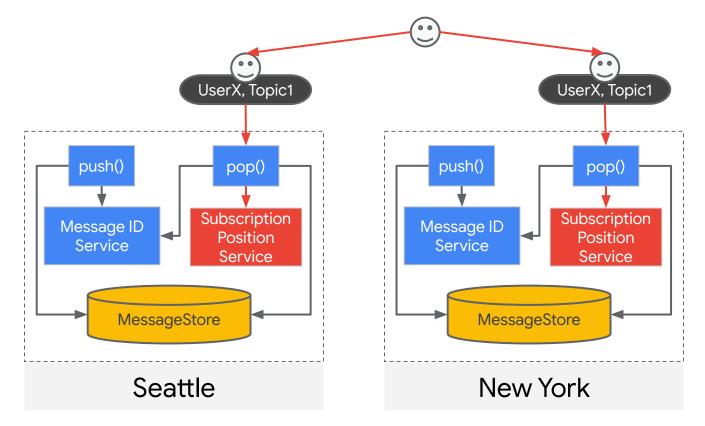
- Distributed components **reliably** and **consistently**:
  - Agree on a single source of truth
  - Identify leaders for specific operations
  - Divide pieces of work
  - Make other decisions
- Unreliable components reliable decisions
- Consistent to decisions, even when sub-components fail
- Recover orphaned datacenters
- Eventual at-most-once semantics
- Paxos, FastPaxos, Raft

### Let's use consensus...

#### **Paxos-based consensus**

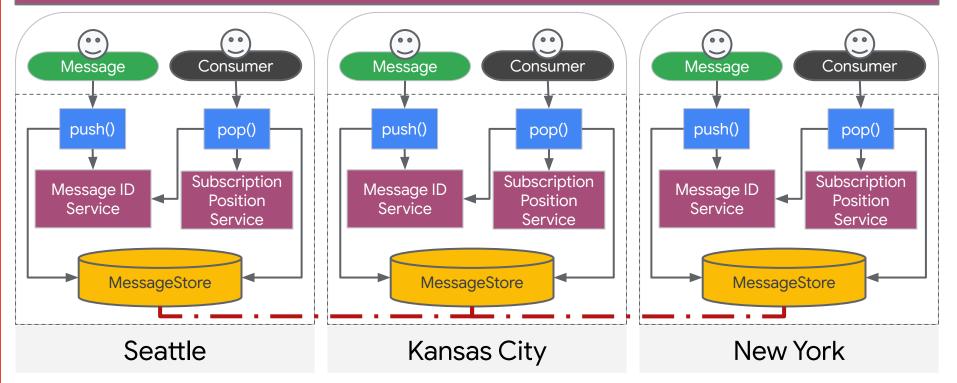


#### Partitioned/Stale Subscription Positions

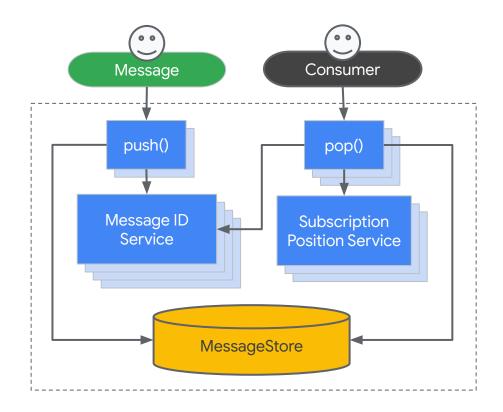


#### Let's use consensus...

#### **Paxos-based consensus**



### **Replicating/Sharding Services**



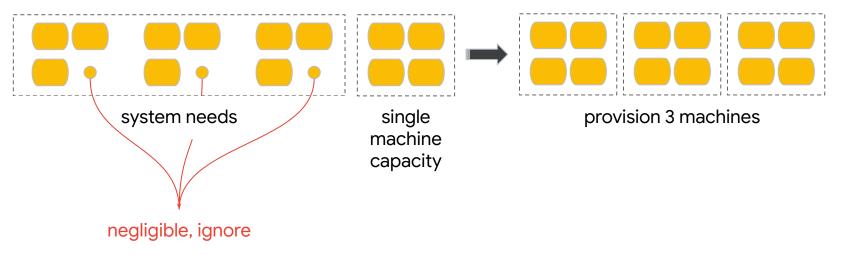
# Breakout Session 3: Provision the System (35 minutes)

# Goal: Identify how many machines you need. Determine if SLOs are viable.

### **Break: 5 Minutes**

## Provisioning

- Provisioning is an art.
- Simplify where possible
- Over-provision by default
- Granularity: units of one machine



# Storage

Message content:

50M msg / day \* 5 kB / msg

= 250 GB / day

IDs:

50M msg / day \* 128 bits / msg **= 800 MB / day** 

Total: ~250 GB / day



Key: Topic ID, Message ID Value: Message Content

Topic ID = 64 bits

Msg ID = 64 bits

Average msg size = 5 kB



# Storage

100 days retention:

250 GB / day \* 100 days

= 25 TB / 100 days

[25 TB / (4 TB HDD / machine)]

= 7 machines

... per DC

... per copy



Key: Topic ID, Message ID Value: Message Content

Topic ID = 64 bits

Msg ID = 64 bits

Average msg size = 5 kB



# Storage

100 days retention:

7 machines / DC / copy

- 7 machines / DC / copy
- \* 2 copies / DC
- \* 3 DCs
- = 42 machines



Key: Topic ID, Message ID Value: Message Content

Topic ID = 64 bits

Msg ID = 64 bits

Average msg size = 5 kB



Machine: 128GB RAM, 2TB SSD 1 x 4TB HDD

#### Which hardware to choose?

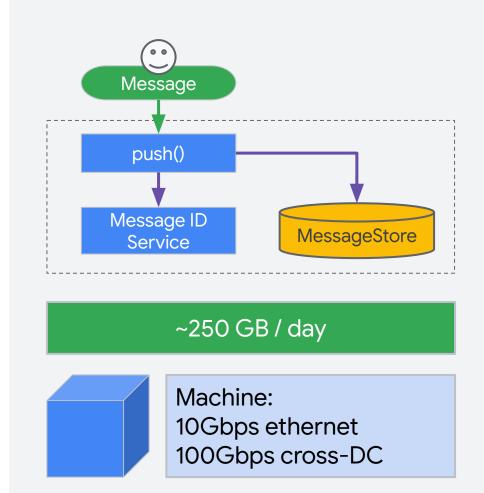


	latency	per-machine	machine count
RAM	0.01ms	128GB	1176
SSD	1ms	2TB	78
HDD	15ms	4TB	42

## Bandwidth: push

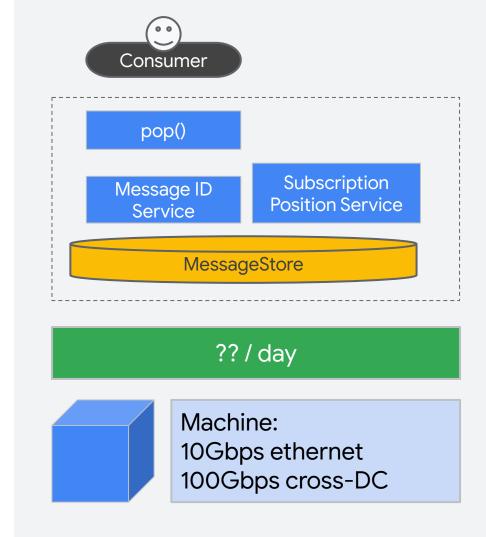
- Peak load = 1.25x avg load
   = 250 GB / day \* 1.25
  - = ~315 GB / day
- 315 GB / day
   = ~4 MB / s
  - = ~30 Mbps inbound
- Outbound ~= Inbound

30 Mbps inbound, 30 Mbps outbound



# Bandwidth: pop

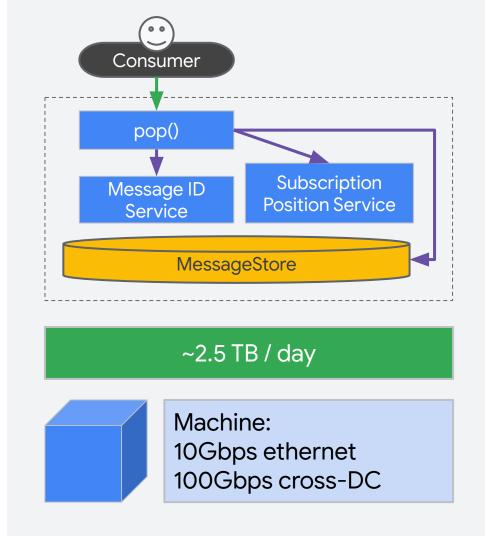
- Avg load
  - = 10k consumers \*
  - 5 topics / consumer \* 10k msg / topic / day \* 5 kB / msg = **2.5 TB / day**



# Bandwidth: pop

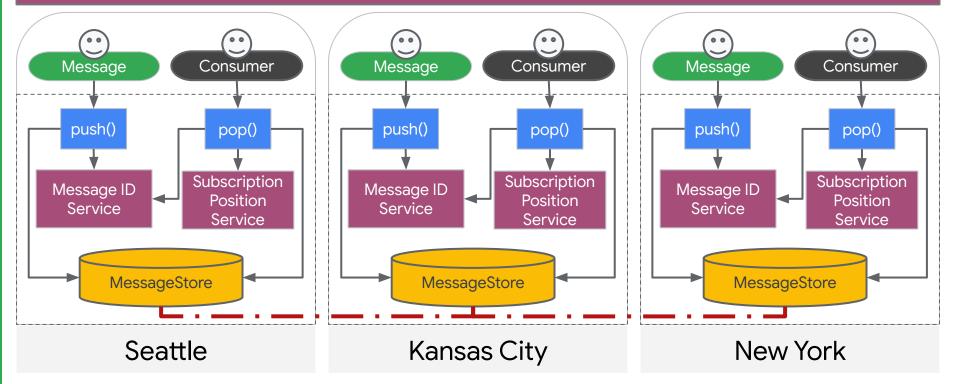
- Peak load = 1.25x avg load
   = 2.5 TB / day \* 1.25
  - = ~3.15 TB / day
- 3.15 TB / day
   = ~37 MB / s
  - = ~300 Mbps outbound
- Internal ~= Outbound

300 Mbps outbound, 300 Mbps internal



#### Is it reliable enough?

#### **Paxos-based consensus**





#### Availability

#### Consistency

(Correctness)

Partition Tolerance

(Latency)

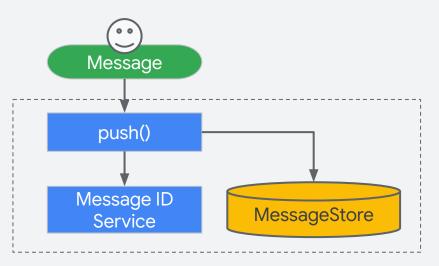
### Latency: push

- Determine ID: ~200ms
- Store message: ~150ms
  - Synchronous
  - Bound by slowest
     connection to remote
     datacenter
- Write message: ~10ms

Total = 200ms + 150ms + 10ms = **360ms** 

#### Reminders:

- 99% ops complete in <500ms
- Paxos takes ~200ms
- Inter-continental = ~150ms
- Local write takes ~10ms



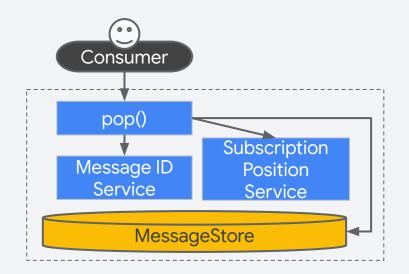
### Latency: pop

- Determine ID: ~0.5ms local,
   ~150ms remote
- Read message: ~15ms local,
   ~150ms remote
- Deliver message: ~negligible
- Update position: ~200ms

Total = 150ms + 150ms + 200ms = **500ms** 

#### Reminders:

- 99% ops complete in <500ms
- Paxos takes ~200ms
- Inter-continental = ~150ms
- Disk seek+read takes ~15ms



# **Bill of Materials** Final count of machines: 2 push +2 pop +3 Message ID Service + 3 Subscription Position Service + 14 MessageStore = 24 per DC \* 3 DCs \* 1.25 (for load spikes) = 90 machines

### Last thoughts

- Start simple and iterate
- See the big picture
- Details, details, details!
- But also, be reasonably pragmatic
- Flexible vs. premature future-proofing
- Cultivate discipline in problem solving approach
- Make data-driven decisions

### Take breaks and enjoy the process!



# **Distributed PubSub**

Non-Abstract Large System Design

